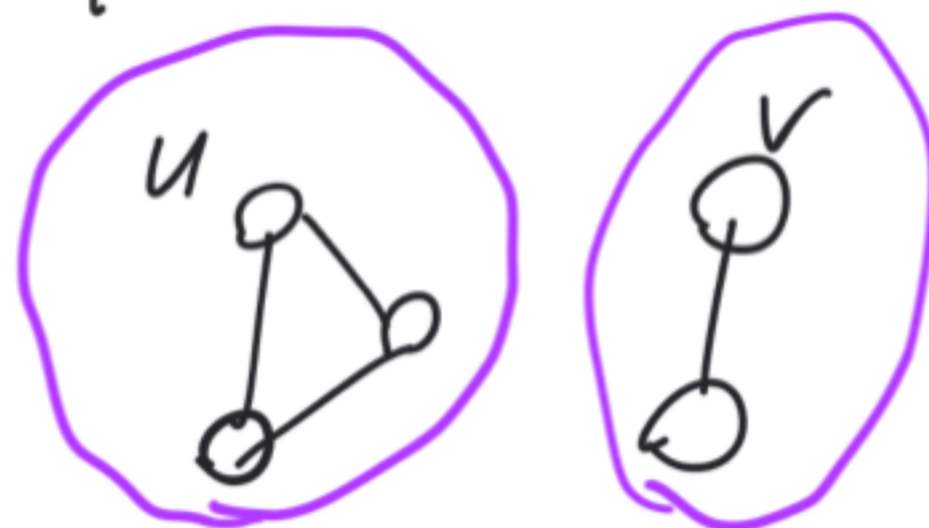
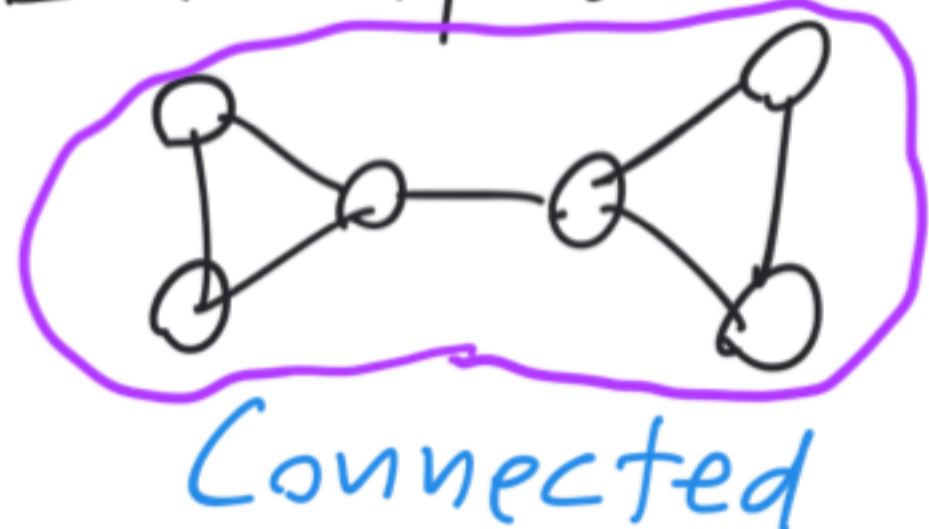


Connectivity and Connected Components

Definition: We say that an undirected graph G is **connected** if for all vertices $u, v \in V(G)$ there exists a path from u to v in G .

Examples:



Connected components

Not connected: There is no path from u to v .

Definition: Given an undirected graph G , a **connected component** of G is a set of vertices $U \subseteq V(G)$ such that:

- 1) $\forall u, u' \in U$ (there is a path from u to u' in G)
- 2) $\forall u \in U \forall v \notin U$ (there is no path from u to v in G)

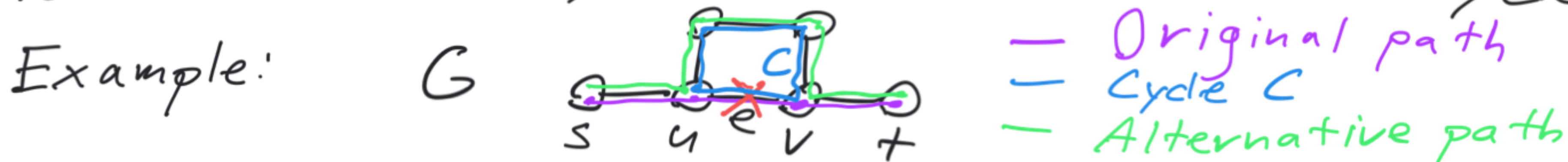
Proposition: Every undirected graph G can be partitioned into connected components.

Connectivity and Connected Components

Q: If we delete an edge $e = \{u, v\}$ from an undirected graph G , how does this affect the connected components of G ?

Claim: If e is contained in a cycle C , deleting e does not affect the connected components of G .

Idea: Instead of using e to go from a vertex s to another vertex t , we can use the rest of the cycle C .



Connectivity and Connected Components

Q: If we delete an edge $e = \{u, v\}$ from an undirected graph G , how does this affect the connected components of G ?

Claim: If $e = \{u, v\}$ is not contained in a cycle then deleting e splits the connected component containing e into two connected components

$U = \{u' \in V(G) : \text{there is a path from } u \text{ to } u' \text{ in } G \text{ which does not contain } e\}$

$V = \{v' \in V(G) : \text{there is a path from } v \text{ to } v' \text{ in } G \text{ which does not contain } e\}$

Observations:

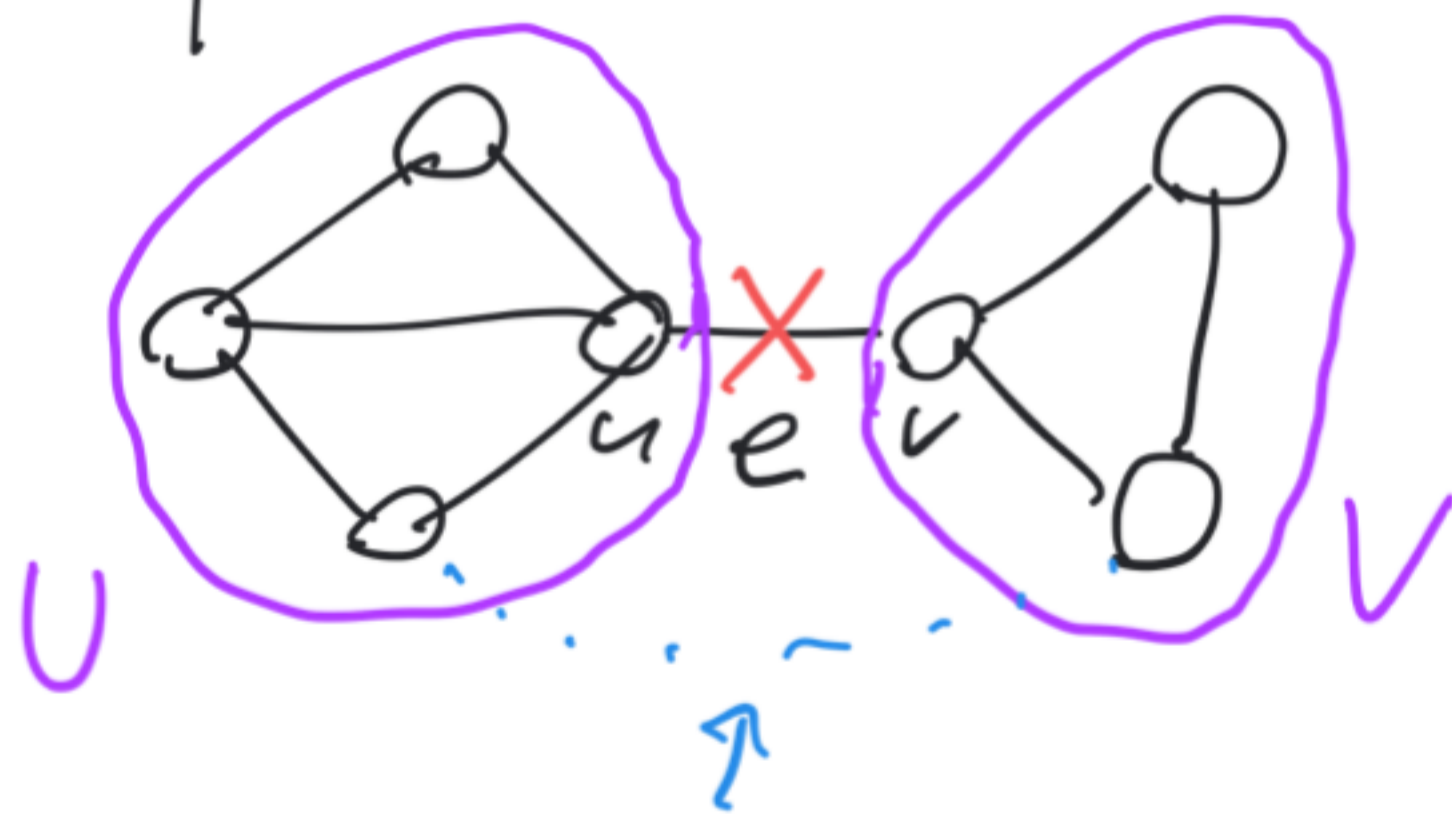
- 1) $U \cup V$ is the original connected component containing e .
- 2) After deleting e , U and V are connected components.
- 3) There is no path from u to v which does not contain e as otherwise e would be contained in a cycle. Thus, $v \notin U$ so U and V are separate connected components.

Connectivity and Connected Components

Q: If we delete an edge $e = \{u, v\}$ from an undirected graph G , how does this affect the connected components of G ?

Claim: If $e = \{u, v\}$ is not contained in a cycle then deleting e splits the connected component containing e into two connected components

Example:



Note: If we had a path from U to V which did not contain e , this would create a cycle containing e .